

CS 421 Lecture 18 – More examples of higher-order functions

- ▶ Combinator programming – “parser combinators”
- ▶ Representing sets as higher-order functions
- ▶ Representing pairs as higher-order functions
- ▶ Building comparators using higher-order functions

Combinator-style programming

Can write complex programs by defining a library of higher-order functions and applying them to one another (and to first-order or built-in functions).

Advantage: easy of creating programs – programs are just expressions

Example: build a parser by writing “parser combinators”.

Parser combinators

Def A parser is a function from token list \rightarrow (token list) option.

Idea is to define functions that build parsers, rather than building parsers "by hand."

E.g. Parser to recognize a single token:

$= \text{fun } s \rightarrow$

let token ~~s~~ = fun cl \rightarrow if cl=[] then None

else if s=hd cl then Some (tl cl)

else None;;

let parsex = token 'x';;

parsex ['x'];; \rightarrow Some []

parsex ['a'];; \rightarrow None

token 'x' ['x']

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let empty = fun lis \rightarrow Some lis

Parser combinators

“Combinators” to combine parsers into larger parsers:

```
let (++) p q = fun cl -> match p cl with None -> None  
                | Some cl' -> q cl';;
```

```
let parsexy = token 'x' ++ token 'y'
```

```
parsexy ['x'; 'y'] → Some []
```

```
parsexy ['x'; 'z'] → None
```

```
parsexy ['y'; 'z'] → None
```

```
parsexy ['x'; 'y'; 'z'] → Some ['z']
```

Parser combinators

```
let (||) p q = fun cl -> match p cl with None -> q cl  
                    | Some cl' -> Some cl';;
```

```
let parsexyorz = parsexy || token 'z'
```

```
parsexyorz ['x'; 'y'] → Some []
```

```
parsexyorz ['z'] → Some []
```

```
parsexyorz ['x', 'z'] → None
```

```
parsexyorz ['z'; 'y'] → Some ['y']
```

Parser combinators

Put this together to define parser for grammar:

$A \rightarrow aB \mid b$

$B \rightarrow cB \mid A$

let rec parseA cl = ((token 'a' ++ parseB) || token 'b') cl

and parseB cl = ((token 'c' ++ parseB) || parseA) cl;;

parseA ['a'; 'c'; 'c'; 'a'; 'b'] \rightarrow Some []

Representing sets as higher-order functions

Def. A set is a function from values to bool.

type intset = int -> bool

E.g. {2} = fun x -> (x=2)

{2,3} = fun x -> (x=2) or (x=3)

Set operations:

(* member: int -> intset -> bool *)

let member n s = s n ;;

(* emptyset: intset *)

let emptyset = fun x -> false

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let
; es = emptyset
;
member 3 es -> false

Representing sets as higher-order functions

(* add: int -> intset -> intset *)

let add n s = fun n' -> s n' or n'=n

(* union: intset -> intset -> intset *)

let union s1 s2 = fun n -> s1 n or s2 n

(* intersection: intset -> intset -> intset *)

let intersection s1 s2 = fun n -> s1 n & s2 n

(* remove: int -> intset -> intset *)

let remove n s = fun n' -> n' ≠ n & s n'

let s3 = add 3 es;;
fun n' -> (es n') | (n'=3)
= n'=3

Representing sets as higher-order functions

(* complement: intset -> intset *)

let complement s = fun n -> not (s n)

(* intsAbove: int -> intset *)

let intsAbove n = fun n' -> (n' > n);;

intsAbove 100 = fun n' -> (n' > 100)

[Note: cannot list elements]

Representing pairs as higher-order functions

Def A *pair* is a value p with a constructor $\text{pair}: \alpha \rightarrow \beta \rightarrow \text{pair}$, and functions $\text{fst}: \text{pair} \rightarrow \alpha$ and $\text{snd}: \text{pair} \rightarrow \beta$ such that $\text{fst}(\text{pair } a \ b) = a$ and $\text{snd}(\text{pair } a \ b) = b$.

let pair a b = fun f → f a b

let fst p = p (fun x → fun y → x)

let snd p = p (fun x → fun y → y)

let pair23 = pair 2 3

⇒ fun f → f 2 3

pair23 : (int → int → int) → int

Building comparators using higher-order functions

$$\text{invert } (\lt) \equiv \gt=$$

Def A *comparator* is a function of type $\alpha * \alpha \rightarrow \text{bool}$.

E.g. (\gt) is a comparator.

$(=)$ is a comparator.

Inverting an ordering:

fun invert ord =

fun (x,y) \rightarrow not (ord (x,y))

or \rightarrow ord (y,x)

Can build specific comparators, e.g.

fun lexorder2 (x,y) (x',y') = x < x' or (x = x' & y < y');;

lexorder2 ('a','b') ('a','c') \rightarrow true

lexorder2 ('a','z') ('b','a') \rightarrow true

lexorder2 ('b','b') ('a','c') \rightarrow false

Building comparators using higher-order functions

But it's more fun to build them using higher-order functions:

```
let or_comp comp1 comp2 = fun (x, y) ->  
  (comp1 (x, y)) or (comp2 (x, y))
```

```
let lte = or_comp (<) (=)
```

```
let and_comp comp1 comp2 = fun (x, y) ->  
  (comp1 (x, y)) & (comp2 (x, y))
```

Building comparators using higher-order functions

```
let lex_comp comp1 comp2 =  
  fun (x,y) (x',y') -> comp1(x,x') or (x=x' & comp2(y y'))
```

```
let lexorder2 = lex_comp (<) (<);;
```

Building comparators using higher-order functions

```
let lex_comp_list comp =
```

```
  let rec aux lis1 lis2 = match (lis1, lis2) with
```

```
    ([], _) -> true
```

```
  | (_, []) -> false
```

```
  | ((x::x'), (y::y')) -> comp x y or (x=y & aux x' y')
```

```
  in aux;;
```

```
let alphalex = lex_comp_list (<);;
```

```
alphalex ['a'; 'b'] ['a'; 'b'; 'c'] → true
```

